

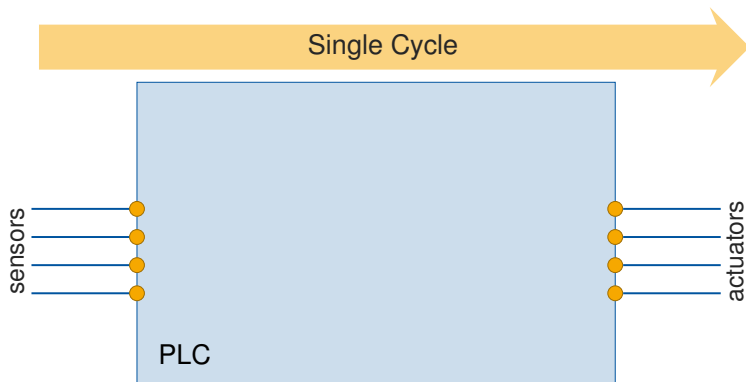


# Verification of Reactive Programs from Industrial Automation

Dimitri Bohlender

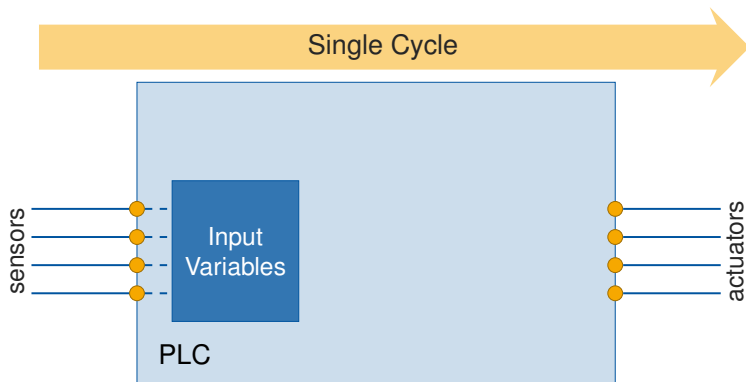
# Programmable Logic Controller (PLC)

- ▶ Tailored to the domain of **industrial automation**
- ▶ Realise **reactive systems**, repeatedly executing the same task



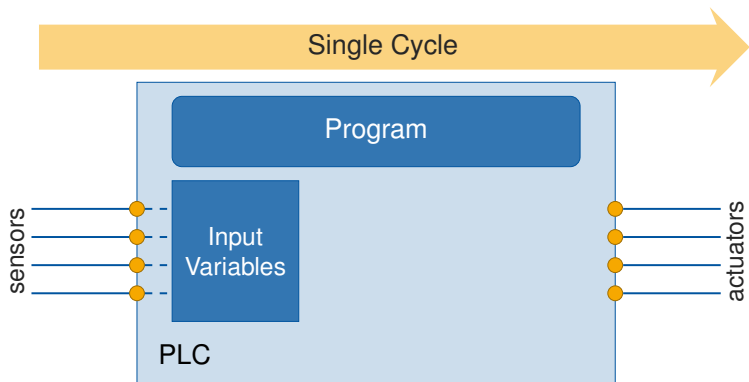
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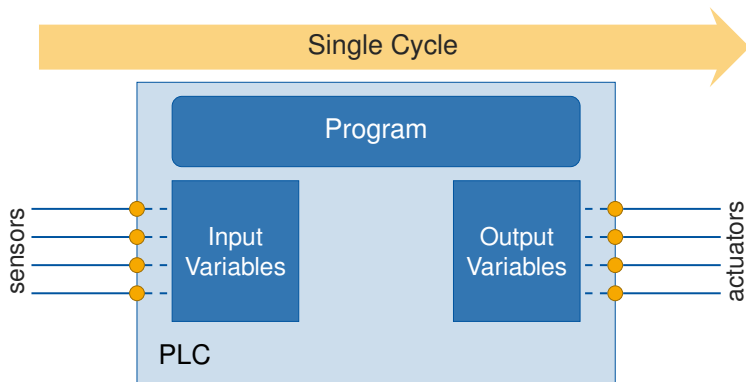
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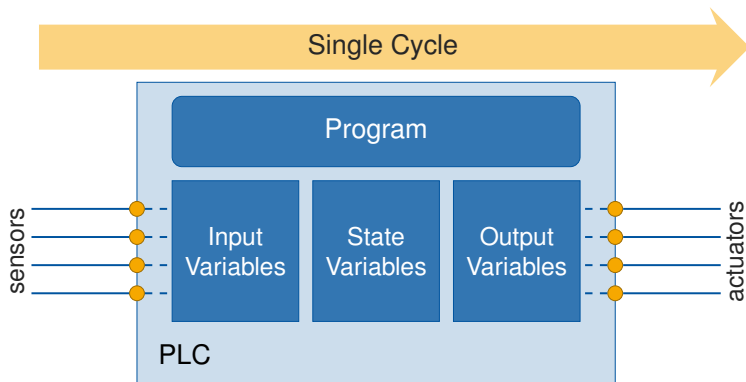
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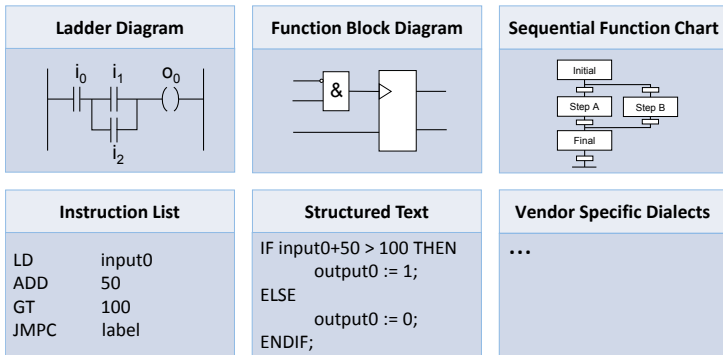
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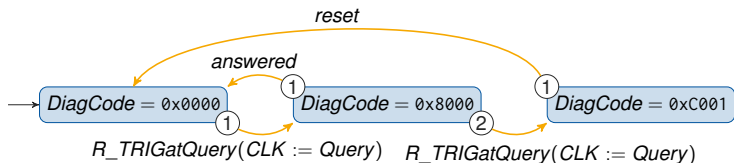
# PLC Software

- ▶ Programming languages standardised in IEC 61131-3
- ▶ Combination of several languages typical



# Verification of Domain-Specific Specifications

- ▶ Specification automata used by the PLCopen



⇒ Characterisation in terms of Constrained Horn-Clauses ✓

- ▶ Analysis of Reset-Behaviour

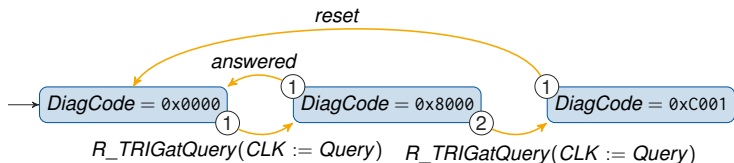
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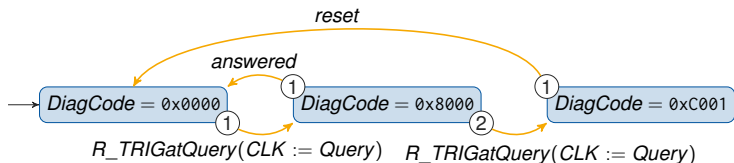
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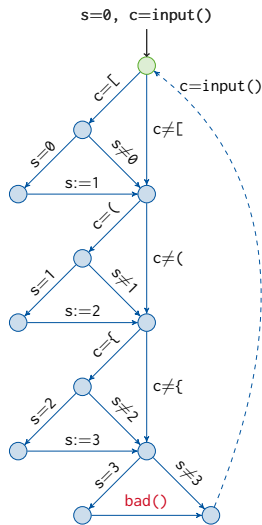
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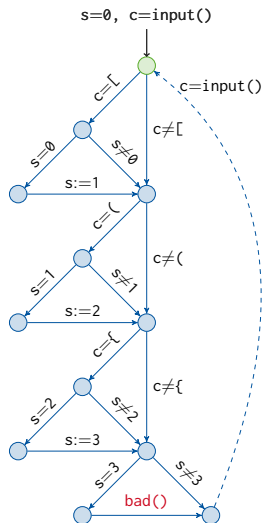
# Exploiting Domain-Specifics in Existing Techniques

- ▶ Consider bug-finding via symbolic execution
- ⇒ CFG-based guidance fails
- ▶ Bad choices **hard to identify** (due to cyclicity)
- ▶ Implicit state machine (over  $s$ )
- ▶ Typical pattern in PLC program modules



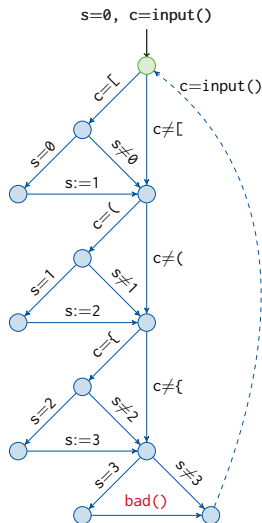
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# Verification of Reactive Programs from Industrial Automation



Dimitri Bohlender, Stefan Kowalewski

## Programmable Logic Controllers (PLCs)

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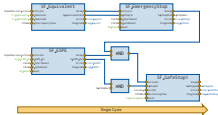


## PLC Software

- Programming languages standardised (IEC 61131-3)
- Combination of several languages typical
- Typically graphical on higher level but textual on lower level

## Observations

- Specifications refer to observable state at cycle-end
- Function blocks exhibit mode-semantics



► Interested?  
Offended?  
⇒ Drop by this poster  
for more details

## PDR-based Model Checking

### PLCopen Automaton

- Specifies safe observable behaviour of a block



- Compliance w.r.t. a transition can often be checked locally
- Encode program in terms of CHCs for a single cycle
- Consider the following (pre-processed) transition



- Check reachability of unsafe behaviour via PDR, where

```

Q1: bool := Q1;
Q2: bool := Q2;
Q3: bool := Q3;
Q4: bool := Q4;
Q5: bool := Q5;
Q6: bool := Q6;
Q7: bool := Q7;
Q8: bool := Q8;
Q9: bool := Q9;
Q10: bool := Q10;
Q11: bool := Q11;
Q12: bool := Q12;
Q13: bool := Q13;
Q14: bool := Q14;
Q15: bool := Q15;
Q16: bool := Q16;
Q17: bool := Q17;
Q18: bool := Q18;
Q19: bool := Q19;
Q20: bool := Q20;
Q21: bool := Q21;
Q22: bool := Q22;
Q23: bool := Q23;
Q24: bool := Q24;
Q25: bool := Q25;
Q26: bool := Q26;
Q27: bool := Q27;
Q28: bool := Q28;
Q29: bool := Q29;
Q30: bool := Q30;
Q31: bool := Q31;
Q32: bool := Q32;
Q33: bool := Q33;
Q34: bool := Q34;
Q35: bool := Q35;
Q36: bool := Q36;
Q37: bool := Q37;
Q38: bool := Q38;
Q39: bool := Q39;
Q40: bool := Q40;
Q41: bool := Q41;
Q42: bool := Q42;
Q43: bool := Q43;
Q44: bool := Q44;
Q45: bool := Q45;
Q46: bool := Q46;
Q47: bool := Q47;
Q48: bool := Q48;
Q49: bool := Q49;
Q50: bool := Q50;
Q51: bool := Q51;
Q52: bool := Q52;
Q53: bool := Q53;
Q54: bool := Q54;
Q55: bool := Q55;
Q56: bool := Q56;
Q57: bool := Q57;
Q58: bool := Q58;
Q59: bool := Q59;
Q60: bool := Q60;
Q61: bool := Q61;
Q62: bool := Q62;
Q63: bool := Q63;
Q64: bool := Q64;
Q65: bool := Q65;
Q66: bool := Q66;
Q67: bool := Q67;
Q68: bool := Q68;
Q69: bool := Q69;
Q70: bool := Q70;
Q71: bool := Q71;
Q72: bool := Q72;
Q73: bool := Q73;
Q74: bool := Q74;
Q75: bool := Q75;
Q76: bool := Q76;
Q77: bool := Q77;
Q78: bool := Q78;
Q79: bool := Q79;
Q80: bool := Q80;
Q81: bool := Q81;
Q82: bool := Q82;
Q83: bool := Q83;
Q84: bool := Q84;
Q85: bool := Q85;
Q86: bool := Q86;
Q87: bool := Q87;
Q88: bool := Q88;
Q89: bool := Q89;
Q90: bool := Q90;
Q91: bool := Q91;
Q92: bool := Q92;
Q93: bool := Q93;
Q94: bool := Q94;
Q95: bool := Q95;
Q96: bool := Q96;
Q97: bool := Q97;
Q98: bool := Q98;
Q99: bool := Q99;
Q100: bool := Q100;
    
```

- Local check may yield spurious counterexamples

```

1 R: TRISTATUS(CSA-Q100);
2 SF (R.TRISTATUS-Q) TANK;
3 DIAGCODE=H0000;
4 END_SF;
5 // FAST_ABORTED;
    
```

- If so, check with closed cycle

## Future Work

- Analysis of restart behaviour: Variables may retain their value after restart/power cut. Starting from these new states no new behaviour shall be observable.
- Mode-oriented PDR: Software-oriented PDR variants partition the transition relation by program locations. An analogous partitioning by modes may help with invariants disjunctive over modes.

## Symbolic Execution

### Guided by Mode-Space

- Consider the right-hand program
- Implicit state-machine (state  $s$ )
- Falls on input sequence "111"
- Bad choices hard to identify (cyclicity)
- CFG-based guidance is local, needs bound and degenerates into random search:



- Mode change cannot be enforced arbitrarily

- Also, some branches are exclusive to certain modes

- Better estimation with mode-space & slicing:

